

Logic Programming

Programming Languages

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Outline

- Predicate Calculus
- Theorem Proving
- Logic Programming
- Case Study: Prolog
- Examples
 - Sudoku
 - N-Queens



Logic Programming

- Expressed in a form of symbolic logic
- Applies logical inferencing to produce results
- **Key insight:** Declarative (instead of Procedural)
 - Specification of results are stated
 - Rather than the procedures which can produce them

Predicate Calculus

Preposition

A logical statement that may or may not be true

- Consists of objects and relationships

Predicate Calculus

Logic founded upon prepositions, variables, constants, and quantifiers

- **Variable** – a symbol that can represent different objects at different times
- **Constant** – a symbol that represents one object
- **Quantifier** – a countable amount (for all, there exists)

Propositions == Compound Terms

- Atomic propositions consist of compound terms
- Compound term describes a relation, but is often expressed as a function (can be written as a table)
- Two parts to a compound term
 - **Functor** – function symbol that names the relationship
 - **Parameters** – ordered list (akin to a tuple)
- Examples:

```
student(jon)
like(beth, macOS)
like(chris, windows)
like(will, linux)
```

Proposition Forms

- Propositions can be stated in two forms:
 - Fact – proposition is assumed to be true
 - Query – truth of proposition is to be determined
- Compound Proposition
 - Have two or more atomic propositions
 - Propositions are connected by operators

Logical Operators

Name	Symbol	Example	Meaning
negation	\neg	$\neg a$	not a
conjunction	\cap	$a \cap b$	a and b
disjunction	\cup	$a \cup b$	a or b
equivalence	\equiv	$a \equiv b$	a is equivalent to b
implication	\supset	$a \supset b$	a implies b
	\subset	$a \subset b$	b implies a

Quantifiers

Name	Example	Meaning
universal	$\forall X.P$	For all X, P is true
existential	$\exists X.P$	There exists a value of X such that P is true

Clausal Form

- We will use a standard form for all propositions
- **Antecedent**
 - Right side
 - What must be true
- **Consequent**
 - Left side
 - What could be true
- Example
 - $B_1 \cup B_2 \cup \dots \cup B_n \subset A_1 \cap A_2 \cap \dots \cap A_m$
 - means if all the As are true, then at least one B is true

Theorem Proving

- Given known axioms and theorems...

We should be able to discover new theorems!

- **Resolution**

- A principle of inference that allows inferred propositions to be computed from given propositions
- Unification
 - Finding values for variables in propositions
- Instantiation
 - Assigning temporary values to variables to allow unification
- After instantiation: if matching fails, we may backtrack

Logic Programming

- Declarative
- Non-Procedural

Programs do not state how to do something!

... Programs state what the result will be.

Logic Programming: Sorting

- Describe the characteristics of a sorted list, rather than the process of rearranging a list

$\text{sort}(\text{old_list}, \text{new_list}) \subset$
 $\text{permute}(\text{old_list}, \text{new_list}) \cap \text{sorted}(\text{new_list})$

$\text{sorted}(\text{list}) \subset$
 $\forall_j \text{ such that } 1 \leq j < n, \text{list}(j) \leq \text{list}(j+1)$



Prolog

Prolog

Predominately used in two fields/areas (origin)

- Natural language processing
- Automated theorem proving

Important Terms:

- *Term* constant, variable, or structure
- *Constant* atom or integer
- *Atom* consists of either:
 - A string of letters, digits, and underscores (starts with a-z)
 - A string of printable ASCII characters delimited by ‘

Terms

Variable

- Any string of letters, digits, or underscores starting with a Capital letter

Instantiation

- Binding of a variable to a concrete value
- May be a temporary binding

Structure

- Represents one atomic proposition
functor(parameter, list)

Facts

Facts are used (in part) to define hypotheses

Known as a “Headless Horn” clause

```
female(amy).  
female(stephanie).  
male(will).  
father(larry, amanda).
```


Rules

Rules are used (in part) to define hypotheses

Known as a “Headed Horn” clause

Right side: *antecedent* (**if** part) – can be a conjunction

Left side: *consequent* (**then** part) – single term

```
ancestor(mary,shelley):- mother(mary,shelley).
```

```
parent(X,Y) :- mother(X,Y).
```

```
parent(X,Y) :- father(X,Y).
```

```
grandparent(X,Z) :- parent(X,Y), parent(Y,Z).
```

Goals

For theorem proving, we may just want to learn or derive something interesting (via proving or disproving)

“Headless Horn” notation:

man(fred)

Can also generalize with variables and propositions

father(X, mike)

female(Y)

Approaches to Solving

- *Matching* is the process of proving a proposition
- Proving a subgoal is called *satisfying* the subgoal
- *Bottom-up resolution, forward chaining*
 - Begin with facts and rules of database and attempt to find sequence that leads to goal
 - Works well with a large set of possibly correct answers
- *Top-down resolution, backward chaining*
 - Begin with goal and attempt to find sequence that leads to set of facts in database
 - Works well with a small set of possibly correct answers
- Prolog implementations use backward chaining

Arithmetic

Integer variables and integer operations are supported

is operator

D **is** B * B - 4 * A * C.

Illegal to do variable reassignment!

Sum **is** Sum + X.

Arithmetic Example

```
speed(ford,100).
speed(chevy,105).
speed(dodge,95).
speed(volvo,80).
time(ford,20).
time(chevy,21).
time(dodge,24).
time(volvo,24).
distance(X,Y) :- speed(X,Speed),
                  time(X,Time),
                  Y is Speed * Time.
```

```
distance(chevy, Chevy_Distance).
```

Lists

- Lists is a sequence of any number of elements
- Elements can be atoms, atomic propositions, or other terms (even other lists!)

[apple, orange, pear, peach]

[] – empty list

[X | Y] – list with head X and tail Y

List Operations - Append

append([], List, List).

append([Head | L1], L2, [Head | Out]) :-
 append(L1, L2, Out).

List Operations - Reverse

```
reverse([], []).
```

```
reverse([Head | Tail], List) :-
```

```
    reverse(Tail, Result),
```

```
    append(Result, [Head], List).
```


List Operations - Member

```
member(Elem, [Elem | _]).  
member(Elem, [_ | List]) :-  
    member(Elem, List).
```



Underscore is an anonymous variable



Deficiencies of Logic Programming

Resolution Order Control

- *The order of attempted matches is non-deterministic and all matches would be attempted concurrently*

The Closed-World Assumption

- *The only knowledge is what is in the database*



The Negation Problem

- *Anything not stated in the database is assumed to be false*
- 



Examples

N-Queens

Problem:

Provided an $N \times N$ chess board, place N queens such that none of them can “take” another (for those with a chess background).

For those without a chess background: place N queens on an $N \times N$ chess board such that there is only one queen per row, one queen per column, and no two queens' difference in rows equals their difference in columns.

N-Queens

<https://swish.swi-prolog.org/example/queens.pl>

Traditional Prolog implementation

- Requires a full board definition
- Iteratively makes all constraints one queen at a time
- Relies on extensive list processing operations

https://swish.swi-prolog.org/example/clpfd_queens.pl

Prolog Implementation relying on CLP(FD) library

- **C**onstraint **L**ogic **P**rogramming over **F**inite **D**omain
- Replaces lists with domains and special operations

N-Queens

```
n_queens(N, Qs) :-
```

```
    length(Qs, N), Qs ins 1..N, safe_queens(Qs).
```

```
safe_queens([]).
```

```
safe_queens([Q|Qs]) :-
```

```
    safe_queens(Qs, Q, 1), safe_queens(Qs).
```

```
safe_queens([], _, _).
```

```
safe_queens([Q|Qs], Q0, D0) :-
```

```
    Q0 #\= Q, abs(Q0 - Q) #\= D0, D1 #= D0 + 1,
```

```
    safe_queens(Qs, Q0, D1).
```

Sudoku

Problem:

Given a 9x9 grid subdivided into 3x3 “houses”, place the values 1 through 9 such that

- *Each row contains each value exactly once*
- *Each column contains each value exactly once*
- *Each house contains each value exactly once*

8								
		3	6					
	7			9		2		
	5				7			
				4	5	7		
			1				3	
		1					6	8
		8	5				1	
	9					4		

Sudoku

https://swish.swi-prolog.org/example/clpfd_sudoku.pl

Prolog Implementation relying on CLP(FD) library

Checking Houses:

```
blocks([], [], []).
```

```
blocks([A,B,C|Bs1],  
       [D,E,F|Bs2],  
       [G,H,I|Bs3]) :-  
    all_distinct([A,B,C,D,E,F,G,H,I]),  
    blocks(Bs1, Bs2, Bs3).
```


Sudoku

https://swish.swi-prolog.org/example/clpfd_sudoku.pl

Prolog Implementation relying on CLP(FD) library

Checking Houses:

```
blocks([], [], []).
```

```
blocks([A,B,C|Bs1],  
       [D,E,F|Bs2],  
       [G,H,I|Bs3]) :-  
    all_distinct([A,B,C,D,E,F,G,H,I]),  
    blocks(Bs1, Bs2, Bs3).
```

Sudoku

Sudoku Solver:

sudoku(Rows) :-

length(Rows,9),

maplist(same_length(Rows),Rows),

append(Rows,Vs), Vs ins 1..9,

maplist(all_distinct,Rows),

transpose(Rows,Columns),

maplist(all_distinct,Columns),

Rows = [A,B,C,D,E,F,G,H,I],

blocks(A,B,C), blocks(D,E,F), blocks(G,H,I).

Sudoku

Problem Definition

```
problem(1, [[_ _ _ _ _ _ _],  
            [_ _ _ _ _ 3, _ 8, 5],  
            [_ _ 1, _ 2, _ _ _],  
  
            [_ _ _ 5, _ 7, _ _ _],  
            [_ _ 4, _ _ _ 1, _ _],  
            [_ 9, _ _ _ _ _],  
  
            [5, _ _ _ _ _ 7, 3],  
            [_ _ 2, _ 1, _ _ _ _],  
            [_ _ _ _ 4, _ _ _ 9]]).
```

Sudoku

Problem Solution

`problem(1, Rows), sudoku(Rows).`

9	8	7	6	5	4	3	2	1
2	4	6	1	7	3	9	8	5
3	5	1	9	2	8	7	4	6
1	2	8	5	3	7	6	9	4
6	3	4	8	9	2	1	5	7
7	9	5	4	6	1	8	3	2
5	1	9	2	8	6	4	7	3
4	7	2	3	1	9	5	6	8
8	6	3	7	4	5	2	1	9