

Bindings and Scope

Programming Languages


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Lecture Outline

- Bindings
 - Type Inferencing
 - Type Binding Examples
 - Lifetime
 - Definition
 - Examples
 - Scope
 - Constructs that create scope
 - Static Scope
 - Dynamic Scope
 - Referencing Environments
- 



Bindings

Type Inferencing

Some languages can infer or deduce the type

```
auto v = 4
```

```
var w = true
```

```
val x = 0
```

```
let y = 3.1415
```

```
const z = "hello"
```

Dynamic Type Binding

The type of a variable can change via assignment

```
lst = [1, 2, 3]  # does not matter  
                # what lst was before
```

- Advantages:
 - Flexibility
 - Ease of use
- Disadvantages:
 - Costly (must do type checking all the time)
 - Type error deduction can be difficult/impossible

Case Study: Python

```
x = 4
```

```
x = 4 + ""
```

```
x = [x]
```

```
x = set(x)
```

```
x = {"x" : x}
```

Case Study: C++

```
auto x = 4;
```

```
auto x = static_cast<float>(x);
```

```
auto x = 4.0f;
```

```
x = 3;
```

```
// another example?
```

Case Study: JavaScript

```
var x = 4
```

```
x = "hello"
```

```
let y = 4
```

```
y = "hello"
```

```
const z = 4
```

```
z = 5
```




Lifetime

Lifetime

- The ***lifetime*** of a variable is the time during which it is bound to a particular memory cell
- Lifetimes can be “created” or “destroyed”
 - **Create**: allocation – getting cell(s) of memory
 - **Destroy**: deallocation – putting cell(s) back
- Four different categories of lifetime
 - Static
 - Stack-dynamic
 - Explicit heap-dynamic
 - Implicit heap-dynamic

Static

```
static int x = 4
```

- Storage bound *prior* to program execution
- Binding *cannot change*
- *Advantages*
 - Direct addressing (immediately available)
 - History sensitive
- *Disadvantages*
 - Inflexible
 - Only one instance permitted

Stack Dynamic

`int x = 4`

- Storage binding created when declaration is “executed” on the computer.
- Only the address can change
- *Advantages*
 - Enables recursion
 - Conserves storage (local)
- *Disadvantages*
 - Allocation/deallocation overhead
 - Not history sensitive
 - Indirect addressing (extra instruction)

Explicit Heap Dynamic `new/delete`

- Allocated and deallocated by explicit directives
- Takes effect during execution
- Referenced only through pointers/references
- *Advantages*
 - Dynamic storage management
- *Disadvantages*
 - Inefficient
 - Unreliable

Implicit Heap Dynamic

=

- Allocated and deallocated by assignment statement
- All arrays/objects in Javascript, Perl, PHP
- *Advantages*
 - Flexibility (generic code)
- *Disadvantages*
 - Inefficient – all attributes are dynamic
 - Loss of error detection



Scope

Scope

- Range of statements over which a variable is visible
- The **scope rules** of a language determine how references to names are associated with variables
- To connect a name reference to a variable, you (or the compiler) must find the declaration
- Static Scope:
 - Based on the **structure** of the program
- Dynamic Scope:
 - Based on the **execution** of the program

Constructs that Create Scope

- Blocks
 - Present in most imperative languages
 - Any statement can be a block of statements
 - Automatically creates a new scope
 - **Shadowing** permitted in C/C++ but disallowed in Java/C#
- Let
 - Present in most functional languages
 - Composed of two parts:
 - Binding names to values
 - Using the names declared in the first part

Scope Categories

- *Local variables*
 - **Declared** within a particular unit of the program
- *Nonlocal variables*
 - **Visible** within a particular unit of the program
 - **Not declared** in that unit of the program
- *Global variables*
 - Visible within **all units** of a program
 - Declared in the outermost scope of the program

Static Scope Rules

Search process:

- Search declarations, first locally
- Then in increasingly larger enclosing scopes, until one is found for the given name

Terms:

- **Static ancestors:** the enclosing static scopes
- **Static parent:** the nearest static ancestor

Variables can be hidden from a unit by having a "closer" variable with the same name (**shadowing**)

Static Scope Example

```
int x = 3;
int main() {
    std::cout << x << '\n';
    int x = 4;
    {
        int x = 5;
        {
            std::cout << x << '\n';
            int x = 6;
            std::cout << x << '\n';
        }
        std::cout << x << '\n';
    }
    std::cout << x << '\n';
}
```

Dynamic Scope Rules

Search process:

- Based on calling sequences of functions
- Search back through the chain of function calls that forced execution to this point

Notes:

- All visible names must be visible to the function called
- “Temporal” in nature – dependent upon what was most recently accessed

Variables **are hidden** from a unit if one of the same name exists in a closer dynamic scope

Dynamic Scope Example

```
function main() {  
    var o = "enter a number: ";  
    print();  
    var x = input();  
    test();  
}
```

```
function test() {  
    if (x % 2 == 0) {  
        even();  
    } else {  
        odd();  
    }  
}
```

```
function even() {  
    o = "even";  
    print()  
}
```

```
function odd() {  
    o = "odd";  
    print()  
}
```

```
function print() {  
    output(o);  
}
```

Static vs. Dynamic Scope

Static

- *Advantages*

- Possible for compiler to detect errors
- Type checking guaranteed

- *Disadvantages*

- Information more difficult to pass

Dynamic

- *Advantages*

- Convenient

- *Disadvantages*

- Impossible(?) to analyze without running
- All variables visible to *subprograms*
- Poor readability

Static vs. Dynamic Scoping

```
function main() {  
    var x = 1;  
    function funA() {  
        var y = x;  
        print(y);  
    }  
    function funB() {  
        var x = 7;  
        funA();  
    }  
    funB();  
}
```

- What does this program print out with*
- *Static Scoping Rules?*
 - *Dynamic Scoping Rules?*

Referencing Environments

- All variables/names visible at a current program unit is known as a **referencing environment**
- We have seen how the referencing environment can vary for static and dynamic scoping rules
- Static-Scoped Languages:
 - Local variables plus all the visible variables in **all enclosing scopes**
- Dynamic-Scoped Languages:
 - Local variables plus all visible variables in **all active subprograms**
 - Active subprogram: when running but not terminated